

A programming algorithm for a product of arithmetical operators  
on M-3.

Order code:

00 ab :	$(b) \Theta (a) \Rightarrow (b), (R)$	$(\Theta ab)$	} $\Theta = +, -, \cdot, \times, \wedge$ (code $\Theta = 0, 1, 2, 3, 6$ )
10 ab :	$(b) \Theta (a) \Rightarrow (R)$	$(\Theta, ab)$	
20 ab :	$(R) \Theta (a) \Rightarrow (b), (R)$	$(\downarrow \Theta ab)$	
30 ab :	$(R) \Theta (a) \Rightarrow (R)$	$(\downarrow \Theta, ab)$	
<del>40 ab :</del>	<del><math>(b) \Theta (a) \Rightarrow (b), (R), paper</math></del>	<del><math>(\Theta \pi ab)</math></del>	
<del>50 ab :</del>	<del><math>(b) \Theta (a) \Rightarrow (R)</math></del>	<del><math>(\Theta   ab)</math></del>	
40 ab :	$(b) \Theta (a) \Rightarrow (b), (R), paper$	$(\Theta \pi ab)$	
50 ab :	$(b) \Theta (a) \Rightarrow (R)$	$(\Theta   ab)$	
60 ab :	$(R) \Theta (a) \Rightarrow (b), (R), paper$	$(\downarrow \Theta \pi ab)$	
70 ab :	$(R) \Theta (a) \Rightarrow (R)$	$(\downarrow \Theta   ab)$	
24 ab	$(R) \Rightarrow (b); \text{next order } (a)$	$(\Pi y ab)$	
34 ab	next order $\begin{cases} (a) \text{ if } (R) \text{ negative} \\ (b) \text{ if } (R) \text{ positive} \end{cases}$	$(\gamma \pi ab)$	
64 ab	$(R) \Rightarrow (b), paper; \text{next order } (a)$	$(\Pi \gamma \pi ab)$	
77 ab	Stop	$(\mathcal{R} ab)$	

Arithmetical terms

Complete bracketing

In fact not required

Hierarchy of elementary operations

Subterm

Non trivial subterm: not a variable

Different occurrences regarded as not identical subterms

equality (formal)

Let us on a stronger equality notion depending on place  
Chain term or ~~chain~~ chain

$$({}^n a_0 \theta_1 a_1) \theta_2 a_2) \theta_3 a_3) \dots) \theta_n a_n$$

~~Chain~~ (n abbreviation for (... ( with n brackets

Order of a chain:  $n, n \neq 0$

Chain operator:  $c \Rightarrow x, c$  a chain

Order same as that of  $c$

Programming of a chain operator:

$$({}^n a_0 \theta_1 a_1) \theta_2 a_2) \theta_3 a_3) \dots) \theta_n a_n) \Rightarrow x$$

$$1 \theta_1 a_1 \quad a_0$$

$$3 \theta_2 a_2$$

$$3 \theta_3 a_3$$

...

$$3 \theta_{n-1} a_{n-1}$$

$$2 \theta_n a_n \quad x$$

In the case  $n=1, (a_0 \theta a_1) \Rightarrow x,$

for  $x \neq a_1, 1 \theta a_1 \quad a_0$

k. 24 kx x;

for  $x = a_1, 0 \theta a_1 \quad a_0$

Valid also for  $\theta_i = 1+1, 1-1, 1:1, |x|, |A|,$  provided

for  $\theta = |A|$  we consider

as  $0 \theta a b$  the order  $5 \Delta a b$

k. 24 kx b

as  $1 \theta a b$  the order  $5 \Delta a b$

as  $2 \theta a b$  the order  $7 \Delta a b$

k. 24 kx b

as  $3 \theta a b$  the order  $7 \Delta a b$

Problem of factorizing arithmetical operators into chain operators

Step by step, have representation of a product of arithmetical operators as a product of chain operators

In the case of an 1 order chain operators, we intend to the identity of right hand side with first component on left hand side:  $(a \ominus b) \Rightarrow a$

Starting bracket: left bracket not preceded by any symbol different from a left bracket

Inner left bracket: not starting

e.g.  $((((x+y)z)-u):((x+y)z)+u)$

A term is a chain if and only if it does not contain any inner left bracket

Inner (non trivial) subterm: beginning with an inner left bracket

e.g.  $((x+y)z)+u$  in the above example

Block of left brackets: no symbol different from a left bracket between

Let us given a ~~chain~~ product

P:  $\frac{t_1 \Rightarrow x_1}{O_1} \frac{t_2 \Rightarrow x_2}{O_2} \dots \frac{t_n \Rightarrow x_n}{O_n}$

of arithmetical operators

Starting, inner left bracket in P: that is one of  $t_1, t_2, \dots, t_n$

Inner subterm of P: that of one of  $t_1, t_2, \dots, t_n$

Maximal subchain of P: not subterm of other subchain

Real equality of subterms of P: besides being formally equal, they do not contain a variable figuring between them on the right hand side of an operator

e.g.  $\underline{(x+y)z}-u \Rightarrow v, \underline{(x+y)z}+u \Rightarrow w$  for real equality

$\underline{(x+y)z}-u \Rightarrow x, \underline{(x+y)z}+u \Rightarrow y$  for formal equality

$\underline{((x+y)z)-u}:v \Rightarrow u, \underline{((x+y)z)+u}:v \Rightarrow v,$

those underlined  $\underline{\quad}$  formally, those underlined  $\underline{\quad}$  really equal

Example 1:

$$((x-y)z+u) : ((x-y)z-u) \Rightarrow u; \quad \underline{\underline{((x+yz)z-v)w-t}} : \underline{\underline{((x+yz)z-u)v+w}} \Rightarrow w;$$

$$\underline{\underline{((x+yz)z-u)v-w}} : \underline{\underline{((x+yz)z-u)v+w}} \Rightarrow v;$$

$$\underline{\underline{((x+yz)z-u)v+w}} : \underline{\underline{((x+yz)z-u)v-w}} \Rightarrow t$$

Steps: 1. Find the left maximal mchain  $c$  of  $P$  ( $\underline{\underline{\quad}}$ )

1,1 Find left left bracket;

1,2 Find the beginning of the block containing it;

1,3 Count the number of its left brackets

1,4 Find the same number of right brackets on the right of it

If  $c$  is on the beginning of  $P$ , then no factorization needed

2. Find the first mchain  $d$  of  $P$  which is really equal to some ~~max~~ mchain of  $c$  ( $d=c$  allowed)

2,1 Store  $c$  in the memory;

2,2 Go to the left; when ~~finding~~ ~~we~~ we find a proper mterm of  $c$ , replace  $c$  by it in the memory;

2,3 When we pass a variable  $x$  standing on the left hand side of  $\Rightarrow$ , store also the maximal mterm of  $c$  not ~~containing~~ containing  $x$ , provided it is non trivial, otherwise do not go further to the left

# Example

$$(a+b+c)^{\frac{1}{2}} ((a+b+c)^{\frac{1}{2}} - a) ((a+b+c)^{\frac{1}{2}} - b) ((a+b+c)^{\frac{1}{2}} - c) \Rightarrow x (=t^2);$$

$$\underbrace{(a+b+c)^{\frac{1}{2}} ((a+b+c)^{\frac{1}{2}} - c)}_{\text{}} : ab \Rightarrow y (=cot^2 \frac{1}{2})$$

$$\underbrace{(a+b+c)^{\frac{1}{2}} ((a+b+c)^{\frac{1}{2}} - a)}_{\text{}} \underbrace{((a+b+c)^{\frac{1}{2}} - b)}_{\text{}} \underbrace{((a+b+c)^{\frac{1}{2}} - c)}_{\text{}} \Rightarrow x; ab \Rightarrow y;$$

$$\underbrace{(a+b+c)^{\frac{1}{2}} ((a+b+c)^{\frac{1}{2}} - c)}_{\text{}} : y \Rightarrow y$$

$$(a+b+c)^{\frac{1}{2}} \Rightarrow s; \underbrace{s(s-a)(s-b)(s-c)}_{\text{}} \Rightarrow x; ab \Rightarrow y; \underbrace{s(s-c)}_{\text{}} : y \Rightarrow y$$

$$(a+b+c)^{\frac{1}{2}} \Rightarrow s; s-c \Rightarrow c; \underbrace{s(s-a)(s-b)c}_{\text{}} \Rightarrow x; ab \Rightarrow y; sc : y \Rightarrow y$$

$$(a+b+c)^{\frac{1}{2}} \Rightarrow s; s-c \Rightarrow c; s-b \Rightarrow x; \underbrace{s(s-a)c}_{\text{}} \Rightarrow x; ab \Rightarrow y; sc : y \Rightarrow y$$

$$(a+b+c)^{\frac{1}{2}} \Rightarrow s; s-c \Rightarrow c; s-b \Rightarrow x; s-a \Rightarrow y; syxc \Rightarrow x; ab \Rightarrow y; sc : y \Rightarrow y$$

0. $1+b$	a	1+	22	21		14. $1 \times b$	a	1x	22	21
1. $3+c$		3+	23	00		15. $2h$	4y	24	16	25
2. $2 \times \frac{1}{2}$	s	2x	26	27		16. $1 \times c$	s	1x	23	27
3. $1-c$	s	1-	23	27		17. $2:y$	y	2:	25	25
4. $245$	c	24	5	23		20. stop		stop		
5. $1-b$	s	1-	22	27		21. a		a		
6. $247$	x	24	7	24		22. b		b		
7. $1-a$	s	1-	21	27		23. c		c		
10. $2411$	y	24	11	25		24. x		x		
11. $1 \times y$	s	1x	25	27		25. y		y		
12. $3 \times x$		3x	24	00		26. $\frac{1}{2}$		$\frac{1}{2}$		
13. $2 \times c$	x	2x	23	24		27. s		s		

Example 1:

$$\frac{((x+y)z-u) : ((x+y)z+u) \Rightarrow u; \wedge x : ((y+z)u-v) \Rightarrow v;}{y : ((y+z)u-v) \Rightarrow w; ((x-y)z+u)v-w \Rightarrow t}$$

$$(x+y)z-u : ((x+y)z+u) \Rightarrow u; (y+z)u \Rightarrow r; x : (r-v) \Rightarrow v;$$
$$y : (r-v) \Rightarrow w; ((x-y)z+u)v-w \Rightarrow t$$

Example 2:

$$((x+y)z+u) : z \Rightarrow y; ((x+y)z-u)v+w : ((x-y)z+u) \Rightarrow z;$$

$$\frac{((x+y)z-v)w-t : ((x+y)z-u)v+w \Rightarrow w;}{((x+y)z-u)v-w : ((x+y)z-u)v+w \Rightarrow v;}$$

$$((x+y)z-u)v-w : ((x+y)z-u)v+w \Rightarrow v;$$

$$\frac{((x+y)z-u)v+w : ((x+y)z-u)v+w \Rightarrow t}{((x+y)z-u)v+w : ((x+y)z-u)v+w \Rightarrow t}$$

Computation of a determinant on M-3.

$$\# \left| \begin{array}{ccc} a_{11} & \dots & a_{1n} \\ \dots & \dots & \dots \\ a_{n1} & \dots & a_{nn} \end{array} \right| = 2^x y, \quad \frac{1}{2} \leq |y| < 1.$$

Method: elimination of all elements ~~on~~ <sup>above</sup> the main diagonal by choosing the ~~appropriate~~ elements with maximal absolute value

Choose  $a_{ij}$  with maximal  $|a_{ij}|$ ; interchange ~~the~~ column  $j$  with column 1. Subtract first column, multiplied by  $\frac{a_{ij}}{a_{11}}$ , from  $j$ th column, however, by dividing the new elements by 2 and raising  $x$  by 1 each time. Supposing, the determinant is already transformed to

$$\left| \begin{array}{cccccccc} a_{11} & 0 & & & 0 & 0 & & 0 \\ a_{21} & a_{22} & a_{23} & & 0 & 0 & & 0 \\ \dots & \dots & \dots & & \dots & \dots & & \dots \\ a_{i-1,1} & a_{i-1,2} & a_{i-1,3} & \dots & a_{i-1,i-1} & 0 & 0 & \dots & 0 \\ a_{i1} & a_{i2} & a_{i3} & \dots & a_{i,i-1} & a_{ii} & a_{i,i+1} & \dots & a_{in} \\ a_{i+1,1} & a_{i+1,2} & a_{i+1,3} & \dots & a_{i+1,i-1} & a_{i+1,i} & a_{i+1,i+1} & \dots & a_{i+1,n} \\ \dots & \dots & \dots & & \dots & \dots & \dots & & \dots \\ a_{n1} & a_{n2} & a_{n3} & \dots & a_{ni-1} & a_{ni} & a_{n,i+1} & \dots & a_{nn} \end{array} \right|,$$

choose  $\max(|a_{i1}|, |a_{i,i+1}|, \dots, |a_{in}|) = |a_{ij}|$ , if  $j \neq i$ , interchange columns  $i$  and  $j$  (first  $i-1$  elements not necessary), if after interchanging  $a_{ii} = 0$ , then take  $x = j = 0$ , otherwise subtract for  $j = i+1, \dots, n$ , column  $i$ , multiplied by  $\frac{a_{ij}}{a_{ii}}$ , from column  $j$ , divided by new column  $j$  by 2 and raise  $x$  by 1. After doing this for  $i = 1, 2, \dots, n-1$ , multiply diagonal elements and normalize  $y$ .

Computation of a determinant on M-3 operator scheme.

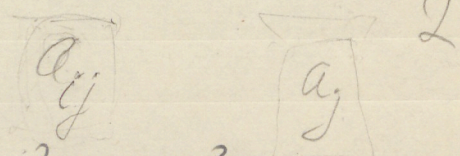
$$0 \Rightarrow x \left\{ \begin{matrix} n-1 \\ i=0 \end{matrix} \right\} \left\{ \begin{matrix} n \\ j=i+1 \end{matrix} \right\} |a_{ii}| < |a_{ij}| \left[ \begin{matrix} n \\ k=i \end{matrix} \right] \left\{ \begin{matrix} a_{ki} \Rightarrow b; a_{kj} \Rightarrow a_{ki} \\ -b \Rightarrow a_{ij} \end{matrix} \right\} \frac{1}{1}$$

$$a_{ii} = 0 \left[ \begin{matrix} 0 \Rightarrow x \\ 0 \Rightarrow y \end{matrix} \right] \frac{1}{2} \left\{ \begin{matrix} n \\ j=i+1 \end{matrix} \right\} \left\{ \begin{matrix} n \\ k=i \end{matrix} \right\} \left\{ \frac{1}{2} a_{kj} - \frac{1}{2} a_{ki} a_{ij} : a_{ii} \right\}$$

~~XXXXXXXXXXXX~~  $x + n^* - i^* \Rightarrow x \left\{ \begin{matrix} a_{ii} \Rightarrow y \\ y a_{ii} \Rightarrow y \end{matrix} \right\} \frac{1}{i=2}$

$$|y| < \frac{1}{2} \left[ \begin{matrix} y : \frac{1}{2} \Rightarrow y \\ x - 1^* \Rightarrow x \end{matrix} \right] \frac{1}{3} \cdot$$

Transformation of



$$0 \Rightarrow x \left\{ \begin{matrix} 1 \Rightarrow y \\ i=0 \end{matrix} \right\} \left\{ \begin{matrix} y a_{ij} \Rightarrow y \\ j=i \end{matrix} \right\} \left\{ \begin{matrix} x+y \Rightarrow x \end{matrix} \right\}$$

$$0 \Rightarrow j$$

1)  $0 \Rightarrow x \left\{ \begin{matrix} 1 \Rightarrow y \\ i=0 \end{matrix} \right\} \left\{ \begin{matrix} i \Rightarrow j \\ \frac{1}{1} y a_{ii} \Rightarrow y \\ j=n \left[ \begin{matrix} j+1 \Rightarrow j \\ \Pi_2^1(\Gamma) \end{matrix} \right] \frac{1}{1} \\ \frac{1}{2} \Pi_2^{-(n-i)}(\Gamma); x+y \Rightarrow x \end{matrix} \right\}$

2)  $0 \Rightarrow x \left\{ \begin{matrix} 1 \Rightarrow y \\ i=0 \end{matrix} \right\} \left\{ \begin{matrix} i \Rightarrow j \\ \frac{1}{1} y a_{ii} \Rightarrow y \\ j=n \left[ \begin{matrix} j+1 \Rightarrow j \\ \Pi_2^1(\Gamma) \end{matrix} \right] \frac{1}{1} \\ \frac{1}{2} B(\Gamma, \Gamma); x+y \Rightarrow x \left[ \begin{matrix} \frac{1}{4} \\ \frac{1}{3} \end{matrix} \right] \frac{1}{4} y a_{ii} \Rightarrow y \left[ \begin{matrix} \frac{1}{4} \\ \frac{1}{4} \end{matrix} \right] \end{matrix} \right\}$

$\frac{3}{\Gamma} \Rightarrow \Gamma$

$$\Gamma \left\{ \begin{matrix} a_{i+1, j} \Rightarrow a_{ij} \\ i=1 \end{matrix} \right\} \Rightarrow \Gamma$$

b\_j^i

